

# Opportunities for reducing communication and improving adaptivity in nonlinear multigrid methods

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This talk: <http://59A2.org/files/20140915-UCDenver.pdf>



# Plan: ruthlessly eliminate communication

- Eliminate, not “aggregate and amortize”

## Why?

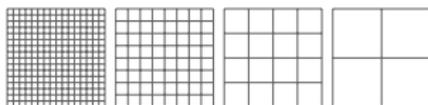
- Local recovery despite global coupling
- Tolerance for high-frequency load imbalance
  - From irregular computation or hardware error correction
- More scope for dynamic load balance

## Requirements

- Must retain optimal convergence with good constants
- Flexible, robust, and debuggable



# Multigrid Preliminaries



**Multigrid** is an  $O(n)$  method for solving algebraic problems by defining a hierarchy of scale. A multigrid method is constructed from:

- 1 a series of discretizations
  - coarser approximations of the original problem
  - constructed algebraically or geometrically
- 2 intergrid transfer operators
  - residual restriction  $I_h^H$  (fine to coarse)
  - state restriction  $\hat{I}_h^H$  (fine to coarse)
  - partial state interpolation  $I_H^h$  (coarse to fine, 'prolongation')
  - state reconstruction  $\mathbb{I}_H^h$  (coarse to fine)
- 3 Smoothers ( $S$ )
  - correct the high frequency error components
  - Richardson, Jacobi, Gauss-Seidel, etc.
  - Gauss-Seidel-Newton or optimization methods



## $\tau$ formulation of Full Approximation Scheme (FAS)

- classical formulation: “coarse grid *accelerates* fine grid” ↘ ↗
- $\tau$  formulation: “fine grid feeds back into coarse grid” ↗ ↘
- To solve  $Nu = f$ , recursively apply

$$\text{pre-smooth} \quad \tilde{u}^h \leftarrow S_{\text{pre}}^h(u_0^h, f^h)$$

$$\text{solve coarse problem for } u^H \quad N^H u^H = \underbrace{I_h^H f^h}_{f^H} + \underbrace{N^H \hat{I}_h^H \tilde{u}^h - I_h^H N^h \tilde{u}^h}_{\tau_h^H}$$

$$\text{correction and post-smooth} \quad u^h \leftarrow S_{\text{post}}^h\left(\tilde{u}^h + I_h^h(u^H - \hat{I}_h^H \tilde{u}^h), f^h\right)$$

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$I_h^H$	residual restriction	$\hat{I}_h^H$	solution restriction
$I_H^h$	solution interpolation	$f^H = I_h^H f^h$	restricted forcing
$\{S_{\text{pre}}^h, S_{\text{post}}^h\}$	smoothing operations on the fine grid		

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- At convergence,  $u^{H*} = \hat{I}_h^H u^{h*}$  solves the  $\tau$ -corrected coarse grid equation  $N^H u^H = f^H + \tau_h^H$ , thus  $\tau_h^H$  is the “fine grid feedback” that makes the coarse grid equation accurate.
- $\tau_h^H$  is *local* and need only be recomputed where it becomes stale.
- Interpretation by Achi Brandt in 1977. many tricks followed



# Model problem: $p$ -Laplacian with slip boundary conditions

- 2-dimensional model problem for power-law fluid cross-section

$$-\nabla \cdot (|\nabla u|^{p-2} \nabla u) - f = 0, \quad 1 \leq p \leq \infty$$

Singular or degenerate when  $\nabla u = 0$

- Regularized variant

$$-\nabla \cdot (\eta \nabla u) - f = 0$$

$$\eta(\gamma) = (\varepsilon^2 + \gamma)^{\frac{p-2}{2}} \quad \gamma(u) = \frac{1}{2} |\nabla u|^2$$

- Friction boundary condition on one side of domain

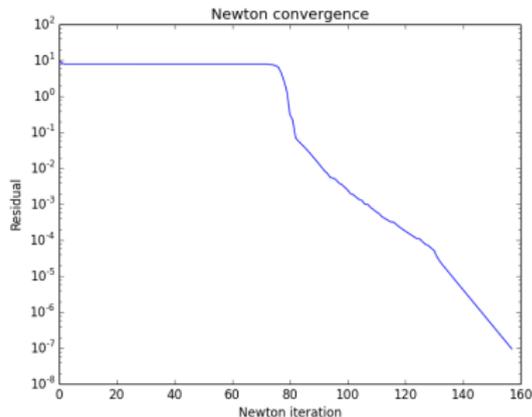
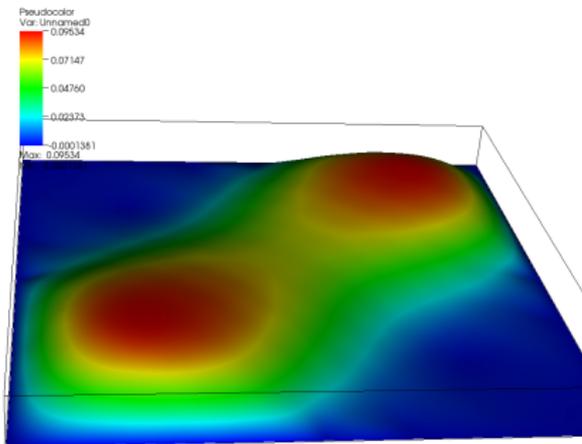
$$\nabla u \cdot n + A(x) |u|^{q-1} u = 0$$



# Model problem: $p$ -Laplacian with slip boundary conditions

- $p = 1.3$  and  $q = 0.2$ , checkerboard coefficients  $\{10^{-2}, 1\}$
- Friction coefficient  $A = 0$  in center, 1 at corners

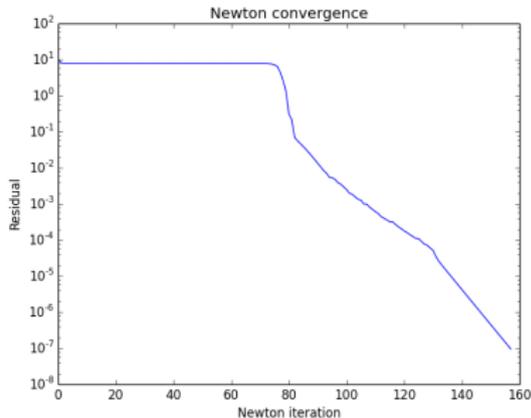
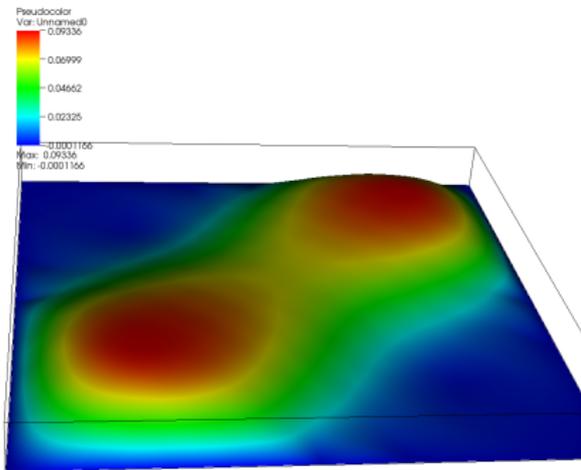
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Cycle: 3



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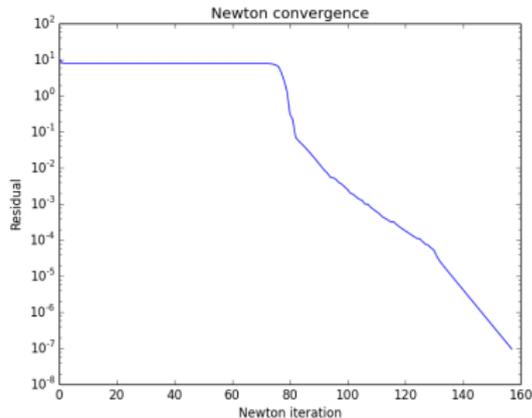
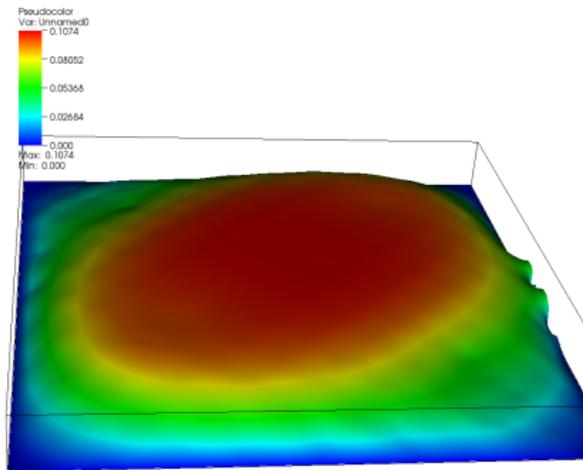
DB: ex15-065.vts  
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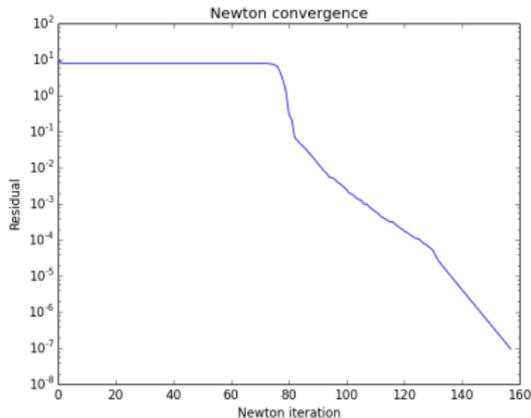
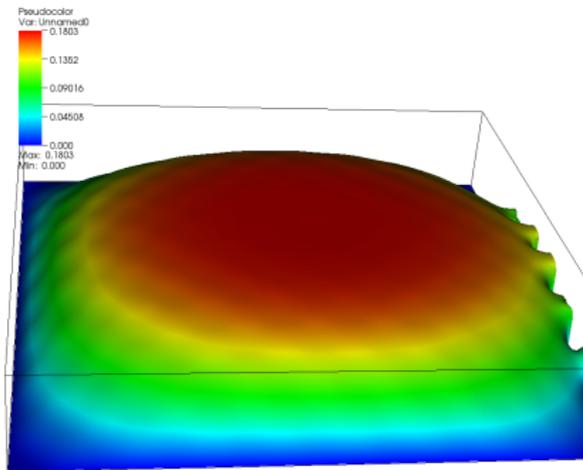
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Cycle: 74



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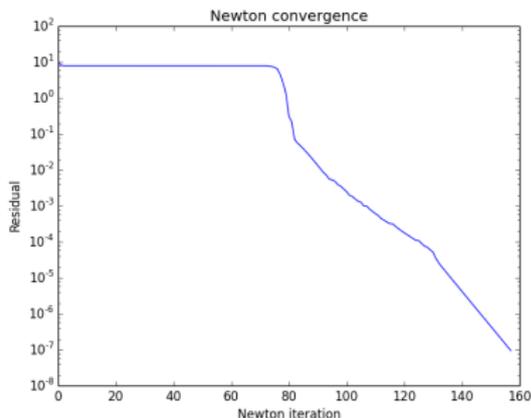
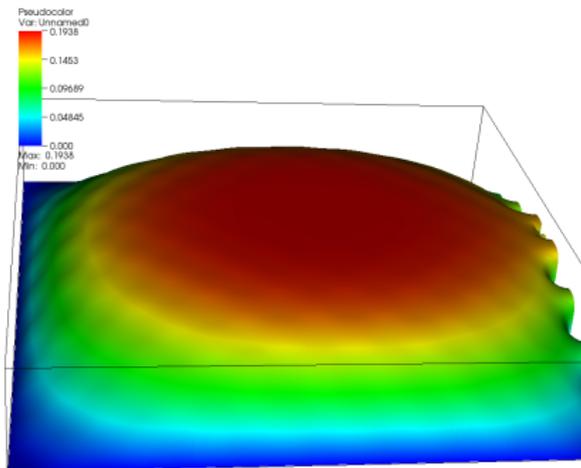
DB: ex15-076.vts  
Cycle: 76



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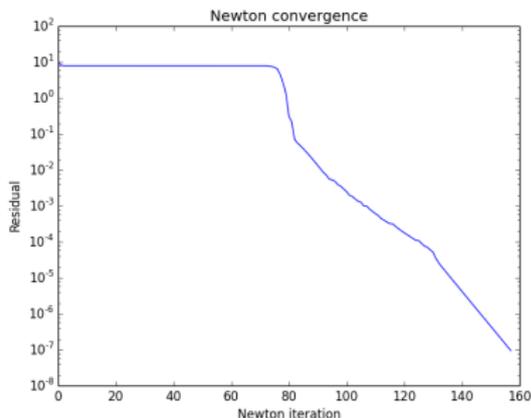
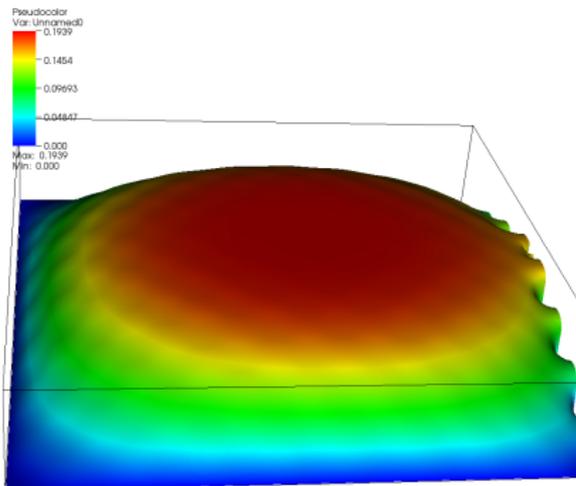
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Cycle: 85



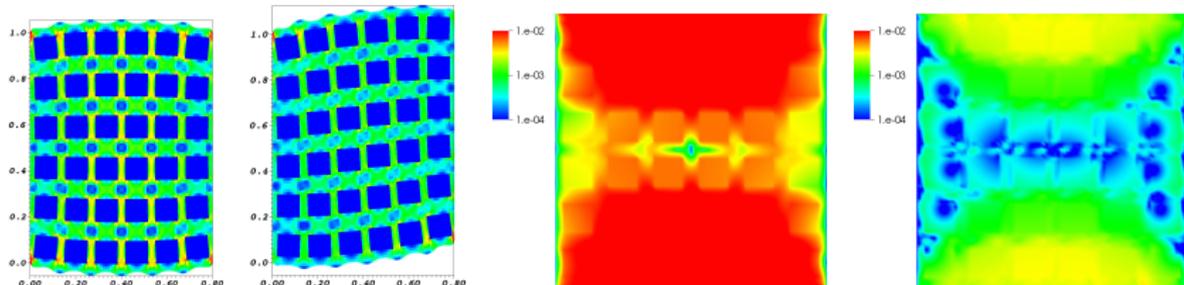
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DB: ex15-115.vts  
Cycle: 115



## $\tau$ corrections



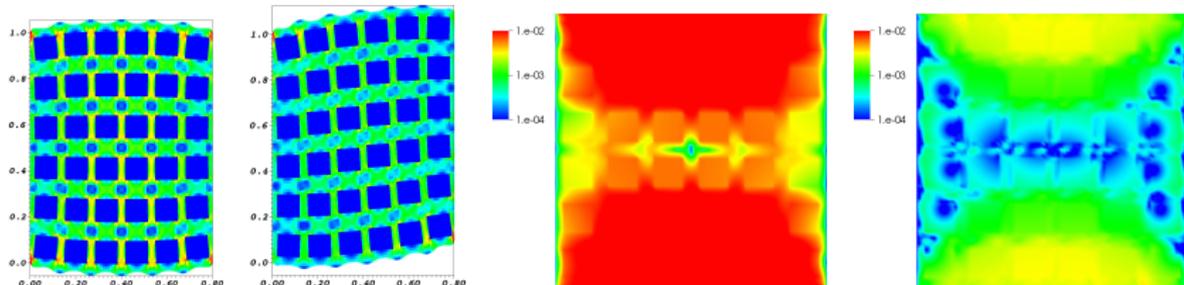
- Plane strain elasticity,  $E = 1000$ ,  $\nu = 0.4$  inclusions in  $E = 1$ ,  $\nu = 0.2$  material, coarsen by  $3^2$ .
- Solve initial problem everywhere and compute  $\tau_h^H = A^H \hat{I}_h^H u^h - I_h^H A^h u^h$
- Change boundary conditions and solve FAS coarse problem

$$N^H \hat{u}^H = \underbrace{I_h^H \hat{f}^h}_{\hat{f}^H} + \underbrace{N^H \hat{I}_h^H \tilde{u}^h - I_h^H N^h \tilde{u}^h}_{\tau_h^H}$$

- Prolong, post-smooth, compute error  $e^h = \hat{u}^h - (N^h)^{-1} \hat{f}^h$
- Coarse grid *with  $\tau$*  is nearly  $10\times$  better accuracy



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# $\tau$ adaptivity: an idea for heterogeneous media

- Applications with localized nonlinearities
  - Subduction, rifting, rupture/fault dynamics
  - Carbon fiber, biological tissues, fracture
- Adaptive methods fail for heterogeneous media
  - Rocks are rough, solutions are not “smooth”
  - Cannot build accurate coarse space without scale separation
- $\tau$  adaptivity
  - Fine-grid work needed everywhere at first
  - Then  $\tau$  becomes accurate in nearly-linear regions
  - Only visit fine grids in “interesting” places: active nonlinearity, drastic change of solution



# Comparison to nonlinear domain decomposition

- ASPIN (Additive Schwarz preconditioned inexact Newton)
  - Cai and Keyes (2003)
  - More local iterations in strongly nonlinear regions
  - Each nonlinear iteration only propagates information locally
  - Many real nonlinearities are activated by long-range forces
    - locking in granular media (gravel, granola)
    - binding in steel fittings, crack propagation
  - Two-stage algorithm has different load balancing
    - Nonlinear subdomain solves
    - Global linear solve
- $\tau$  adaptivity
  - Minimum effort to communicate long-range information
  - Nonlinearity sees effects as accurate as with global fine-grid feedback
  - Fine-grid work always proportional to “interesting” changes

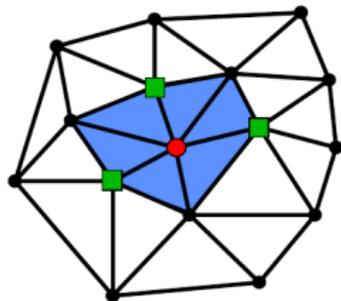


# Nonlinear and matrix-free smoothing

- matrix-based smoothers require global linearization
- nonlinearity often more efficiently resolved locally
- nonlinear additive or multiplicative Schwarz
- nonlinear/matrix-free is good if

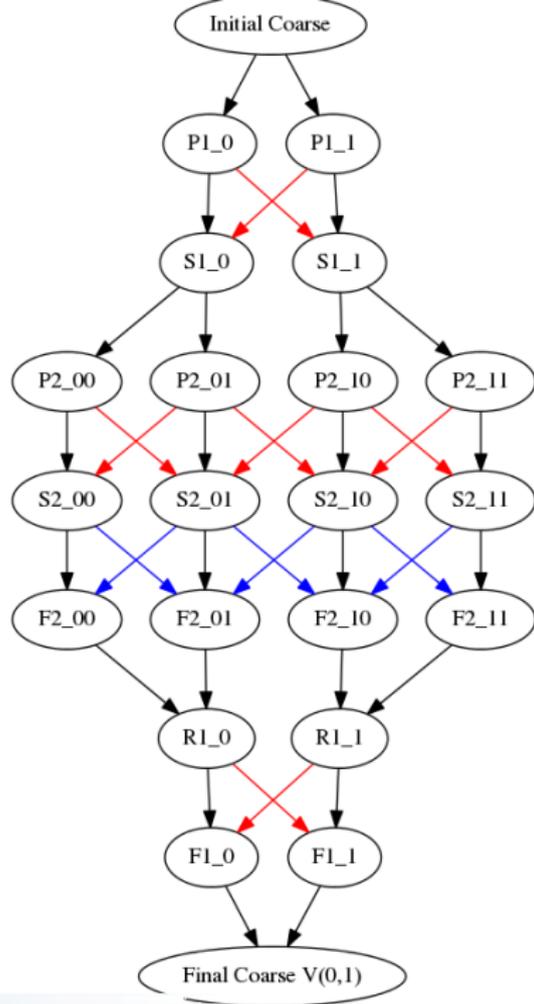
$$C = \frac{(\text{cost to evaluate residual at one "point"}) \cdot N}{(\text{cost of global residual})} \sim 1$$

- finite difference:  $C < 2$
- finite volume:  $C \sim 2$ , depends on reconstruction
- finite element:  $C \sim$  number of vertices per cell
- larger block smoothers help reduce  $C$
- additive correction (Jacobi/Chebyshev/multi-stage)
  - global evaluation, as good as  $C = 1$
  - but, need to assemble corrector/scaling
  - need spectral estimates or wave speeds



# Low communication MG

- **red arrows** can be removed by  $\tau$ -FAS with overlap
- **blue arrows** can also be removed, but then algebraic convergence stalls when discretization error is reached
- no simple way to check that discretization error is obtained
- if fine grid state is not stored, use compatible relaxation to complete prolongation  $P$
- “Segmental refinement” by Achi Brandt (1977)
- 2-process case by Brandt and Diskin (1994)

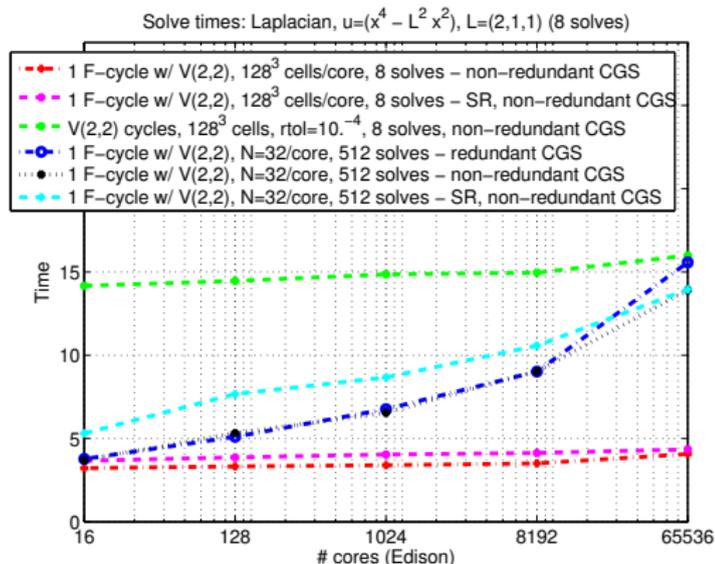


# Segmental refinement: no horizontal communication

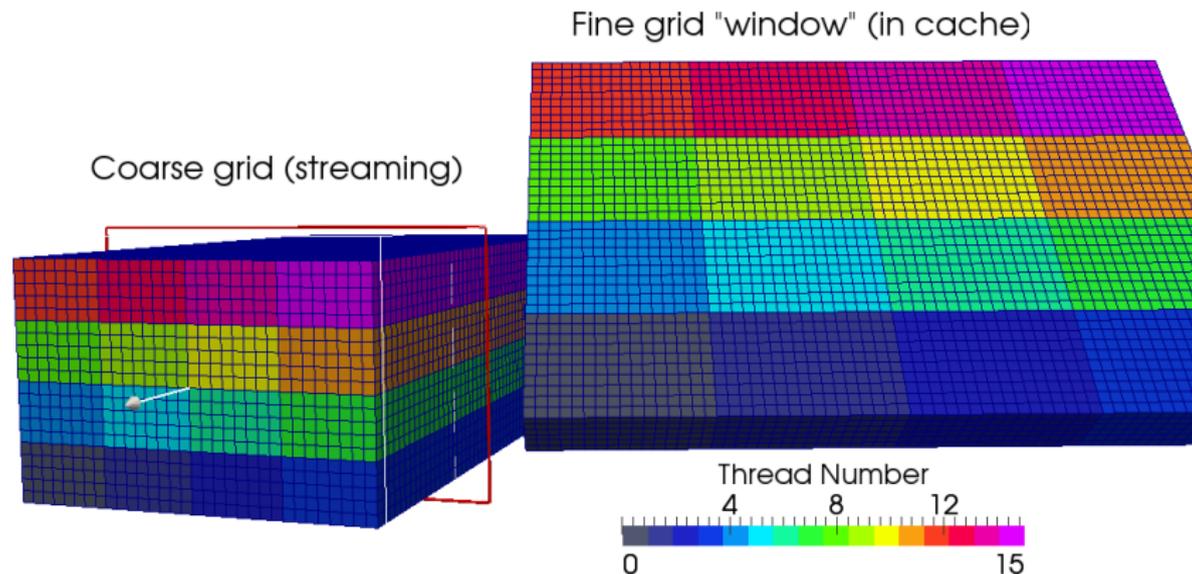
- 27-point second-order stencil, manufactured analytic solution
- 5 SR levels:  $16^3$  cells/process local coarse grid
- Overlap = Base +  $(L - \ell)$  Increment
  - Implementation requires even number of cells—round down.
- FMG with  $V(2,2)$  cycles

Table:  $\|e_{SR}\|_{\infty} / \|e_{FMG}\|_{\infty}$

Increment	Base		
	1	2	3
1	1.59	2.34	1.00
2	1.00	1.00	1.00
3	1.00	1.00 <td 1.00	



# Reducing memory bandwidth



- Sweep through “coarse” grid with moving window
- Zoom in on new slab, construct fine grid “window” in-cache
- Interpolate to new fine grid, apply pipelined smoother ( $s$ -step)
- Compute residual, accumulate restriction of state and residual into coarse grid, expire slab from window



## Arithmetic intensity of sweeping visit

- Assume 3D cell-centered, 7-point stencil
- 14 flops/cell for second order interpolation
- $\geq 15$  flops/cell for fine-grid residual or point smoother
- 2 flops/cell to enforce coarse-grid compatibility
- 2 flops/cell for plane restriction
- assume coarse grid points are reused in cache
- Fused visit reads  $u^H$  and writes  $\hat{I}_h^H u^h$  and  $I_h^H r^h$
- Arithmetic Intensity

$$\frac{\underbrace{15}_{\text{interp}} + \underbrace{2 \cdot (15 + 2)}_{\text{compatible relaxation}} + \underbrace{2 \cdot 15}_{\text{smooth}} + \underbrace{15}_{\text{residual}} + \underbrace{2}_{\text{restrict}}}{3 \cdot \text{sizeof}(\text{scalar}) / \underbrace{2^3}_{\text{coarsening}}} \gtrsim 30 \quad (1)$$

- Still  $\gtrsim 10$  with non-compressible fine-grid forcing



## Regularity

Accuracy of recovery depends on operator regularity

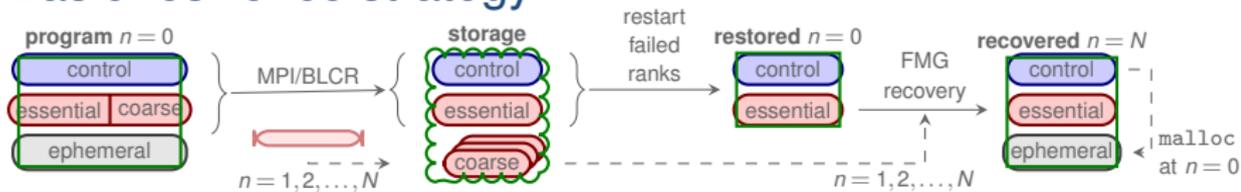
- Even with regularity, we can only converge up to discretization error, unless we add a *consistent* fine-grid residual evaluation
- Visit fine grid with some overlap, but patches do not agree exactly in overlap
- Need decay length for high-frequency error components (those that restrict to zero) that is bounded with respect to grid size
- Required overlap  $J$  is proportional to the number of cells to cover decay length
- Can enrich coarse space along boundary, but causes loss of coarse-grid sparsity
- Brandt and Diskin (1994) has two-grid LFA showing  $J \lesssim 2$  is sufficient for Laplacian
- With  $L$  levels, overlap  $J(k)$  on level  $k$ ,

$$2J(k) \geq s(L - k + 1)$$

where  $s$  is the smoothness order of the solution or the discretization order (whichever is smaller)



# Basic resilience strategy



control contains program stack, solver configuration, etc.

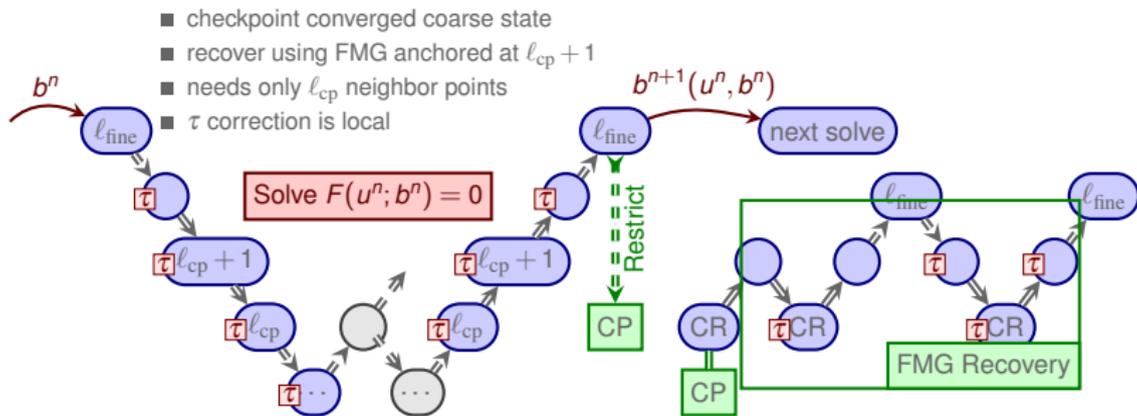
essential program state that cannot be easily reconstructed:  
time-dependent solution, current optimization/bifurcation  
iterate

ephemeral easily recovered structures: assembled matrices,  
preconditioners, residuals, Runge-Kutta stage solutions

- Essential state at time/optimization step  $n$  is **inherently globally coupled** to step  $n - 1$  (otherwise we could use an explicit method)
- *Coarse* level checkpoints are orders of magnitude smaller, but allow rapid recovery of essential state
- FMG recovery needs only **nearest neighbors**



# Multiscale compression and recovery using $\tau$ form



- Normal multigrid cycles visit all levels moving from  $n \rightarrow n + 1$
- FMG recovery only accesses levels finer than  $l_{CP}$
- Lightweight checkpointing for transient adjoint computation
- Postprocessing applications, e.g., in-situ visualization at high temporal resolution in part of the domain



## First-order cost model for FAS resilience

Extend first-order locality-unaware model of Young (1974):

$t_W$  time to write a heavy fine-grid checkpointed state

$t_R$  time to read back lost state

$R$  fraction of forward simulation needed for recomputation from a saved state

$P$  the heavy checkpoint interval

$M$  mean time to failure

Neglect cost of I/O for lightweight coarse-grid checkpoints

$$\text{Overhead} = 1 - \text{AppUtilization} = \underbrace{\frac{t_W}{P}}_{\text{writing}} + \underbrace{\frac{t_R}{M}}_{\text{reading after failure}} + \underbrace{\frac{RP}{2M}}_{\text{recomputation}}$$

Minimized for a heavy checkpointing interval  $P = \sqrt{2Mt_W/R}$

$$\text{Overhead}^* = \sqrt{2t_W R/M} + t_R/M$$

where the first term is always larger than the second. Conventional checkpointing schemes store only fine-grid state, thus  $R = 1$  (recovery costs the same as initial computation).



## Other uses

- Transient adjoints
  - Adjoint model runs backward-in-time, needs state from solution of forward model
  - Status quo: hierarchical checkpointing
  - Memory-constrained and requires computing forward model multiple times
  - If forward model is stiff, each step has global dependence
  - Compression via  $\tau$ -FAS accelerates recomputation, can be local
- Visualization and analysis
  - Targeted visualization in small part of domain
  - Interesting features emergent so can't predict where to look



# Outlook on $\tau$ -FAS adaptivity and compression

- Benefits of AMR without fine-scale smoothness
- Coarse-centric restructuring is a major interface change
- Nonlinear smoothers (and discretizations)
  - Smooth in neighborhood of “interesting” fine-scale features
  - Which discretizations can provide efficient matrix-free smoothers?
  - Does there exist an efficient smoother based on element Neumann problems?
- Dynamic load balancing
- Reliability of error estimates for refreshing  $\tau$ 
  - We want a coarse indicator for whether  $\tau$  needs to change
- Worthwhile for resilience and to better use hardware

